

Towards a general construction of playgrounds

Clovis Eberhart

École Normale Supérieure de Cachan

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Outline

- 1 Introduction
- 2 Building plays
- 3 A correctness criterion
- 4 Conclusion

Introduction

- notion of playground

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- instances with similar constructions

Introduction

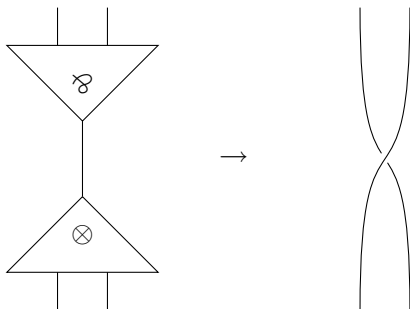
- notion of playground
- instances with similar constructions
- general and automated construction of playgrounds

Introduction

- notion of playground
- instances with similar constructions
- general and automated construction of playgrounds
- simple playground to demonstrate the categorical combinatorics methods

An MLL interaction net

Reduction in MLL interaction nets



Let us take a closer look at the interaction net

The base category

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- objects of dimension 0 are called *channels*
- objects of dimension 1 are called *players*
- objects of dimension at least 2 are called *moves*

The base category

In the case of MLL interaction nets:

*

dimension 0

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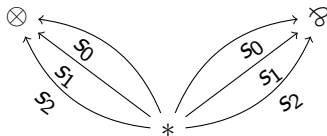
dimension 1

*

dimension 0

The base category

In the case of MLL interaction nets:



dimension 1

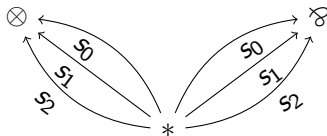
dimension 0

The base category

In the case of MLL interaction nets:

cut

dimension 2

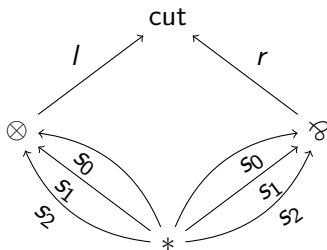


dimension 1

dimension 0

The base category

In the case of MLL interaction nets:



dimension 2

dimension 1

dimension 0

with the relations $l \circ s_j = r \circ s_j$.

Positions as presheaves

A presheaf on a category \mathcal{C} is a functor $F : \mathcal{C}^{op} \rightarrow \text{Set}$.
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In the case of MLL interaction nets, positions are presheaves of dimension 1 such that there are never more than two morphisms from players to a given channel.

Positions as presheaves

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- objects: pairs (c, x) where c is an object of \mathcal{C} and $x \in F(c)$
- morphisms: $u : (c, x) \rightarrow (c', x')$ if $u : c \rightarrow c'$ is a morphism of \mathcal{C} such that $F(u)(x') = x$.

Category of elements

Seeds

We assume given, for every move m of \mathcal{C} , a cospan in the category of presheaves $Y \rightarrow M \leftarrow X$ (called a “seed”), where X and Y are positions, M is the representable presheaf corresponding to m and that verifies the following properties:

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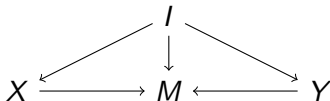
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- $X(p) + Y(p) \rightarrow M(p)$ is surjective for every player p
- it has a “canonical interface” I of dimension 0 such that:

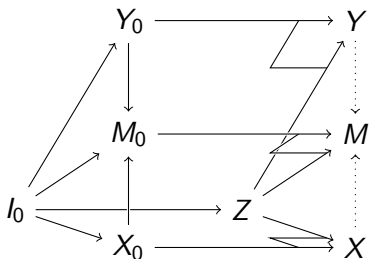


commutes.

Seeds

Moves

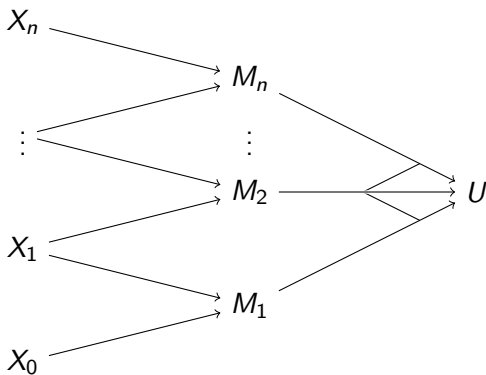
A move is a cospan $Y \rightarrow M \leftarrow X$ built from a seed $Y_0 \rightarrow M_0 \leftarrow X_0$ with canonical interface I by pushing out along some “suitable” position Z in the following way:



Moves

Building plays

A play is a composition of moves in the bicategory of cospans.



A first example

A second example

Problem

- plays are cospans $Y \rightarrow U \leftarrow X$, but when is such a cospan a play?

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- plays are cospans $Y \rightarrow U \leftarrow X$, but when is such a cospan a play?
- need for a correctness criterion.

Past

For every seed $Y \xrightarrow{s} M \xleftarrow{t} X$, we define the past of M to be the set:

$$\text{past}(M) = \bigcup_{p \in \mathcal{C}} M(p) - Y(p)$$

Cores, core separation

A core of a presheaf U is a move μ of $\text{el}(U)$ such that, if $f : \mu \rightarrow x$ in $\text{el}(U)$, then $x = \mu$ and $f = \text{id}_\mu$.

A presheaf $U \in \widehat{\mathcal{C}}^f$ is core-separating if for all cores $\mu \neq \mu'$ in $\text{el}(U)$, the pullback of μ along μ' is a position.

Local 1-injectivity

A presheaf U on \mathcal{C} is locally 1-injective iff for every seed $Y \xrightarrow{s} M \xleftarrow{t} X$ with canonical interface $u : I \rightarrow M$ and for all corresponding core $\mu \in \text{el}(U)$ (seen as a morphism $\mu : M \rightarrow U$), if $x \neq y \in M$ are such that $\mu(x) = \mu(y)$, then x, y are in the image of u .

Partitioning players and channels

For every seed, we partition players in the following way:

- consumed players: $\text{Co}(M)(p) = X(p) \setminus Y(p)$
- created players: $\text{Cr}(M)(p) = Y(p) \setminus X(p)$
- surviving players: $\text{Sr}(M)(p) = X(p) \cap Y(p)$

We do the same for channels.

Initial and final players and channels

For every presheaf $U \in \widehat{\mathcal{C}}^f$, we define the set of its initial and final players by:

- $\text{Init}(U)(p) = \{x \in U(p) \mid \nexists m \in \mathcal{C}, \tilde{m} \in U(m), x \in \text{Cr}(\tilde{m})(p)\}$
- $\text{Fin}(U)(p) = \{x \in U(p) \mid \nexists m \in \mathcal{C}, \tilde{m} \in U(m), x \in \text{Co}(\tilde{m})(p)\}$

The causal graph

We define the causal graph G_U that has:

G_U is source-linear if for all $x \rightarrow \mu, x \rightarrow \mu', \mu = \mu'$.

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 - $x \rightarrow x \cdot s$ for every player x and $s : p \rightarrow *$

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 - $\mu \rightarrow x$ for every core μ and $x \in \text{Co}(\mu)(*) \cup \text{Sr}(\mu)(*)$
 - $\mu \rightarrow \mu'$ for all cores $\mu \neq \mu'$ when there is a player x in $\text{Co}(\mu) \cap (\text{Co}(\mu') \cup \text{Sr}(\mu'))$

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Assumed

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- for every seed M , $\text{past}(M)$ only contains moves and players
- there is no isolated channel
- every seed $Y \xrightarrow{s} M \xleftarrow{t} X$ has a canonical interface $I = X_0$
($X_0(x) = X(x)$ in dimension 0, $X_0(x) = \emptyset$ otherwise)

The correctness criterion

Correctness Criterion

A cospan $Y \hookrightarrow U \leftarrow X$ is a play iff the following conditions are met:

- U is core-separating and locally 1-injective
- X contains exactly the initial players and channels of U
- Y contains exactly the final players and channels of U
- G_U is source-linear and acyclic

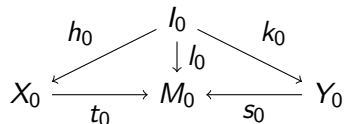
Lemma

Lemma

Let U be a presheaf on \mathcal{C} and μ be maximal in G_U (i.e., there is no path from μ to any other core). Assume that U is core-separating. Then for all $c \in \text{el}(U)$, $U(c) - \text{past}(\mu) = (U \setminus \mu)(c)$.

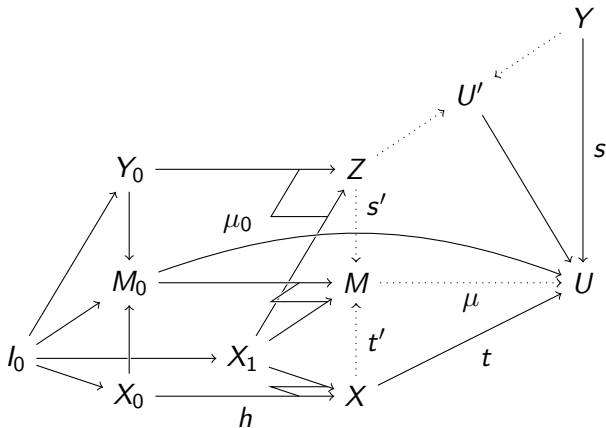
Sketch of the proof

Take μ_0 maximal in G_U , and be



its canonical interface.

Sketch of the proof



Conclusion

- MLL plays as cospans of presheaves

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- garbage collection